Quantum Recursion and Second Quantisation

Mingsheng Ying

State Key Laboratory of Computer Science

Outline

- 1. Introduction
- 2. Recursive Programming
- 3. Classical Recursion in Quantum Programming
- 4. Quantum Control Flow
- 5. Motivating Example: Recursive Quantum Walks
- 6. Second Quantisation
- 7. Semantics of Quantum Recursion
- 8. Conclusion

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IBM, Google, Intel, Microsoft building quantum computers

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- ► IBM Q: 5 quantum bits (qubits)
- Google: quantum supremacy

Will you be quantum Alan Turing?

Model of quantum computation — Quantum Turing machine

- [1] P. Benioff, The computer as a physical system: A microscopic quantum mechanical Hamiltonian model of computers as represented by Turing machines, *J. of Statistical Physics* 1980.
- [2] I. Yu. Manin, Computable and Noncomputable (in Russian), Sov. Radio 1980.
- [3] R. Feynman, Simulating physics with computers, *Int. J. of Theoretical Physics* 1982.
- [4] D. Deutsch, Quantum theory, the Church-Turing principle and the universal quantum computer, *Proc. of the Royal Society of London A* 1985.

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Mathematical Logic: Recursion

A long history in Mathematics!

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Recursive programming

Put forward and implement the recursive procedure as an ALGOL60 language construct

[5] E. W. Dijkstra, Recursive programming, *Numerische Mathematik* 1960.

[6] E. G. Daylight, Dijkstra's rallying cry for generalization: The advent of the recursive procedure, late 1950s - early 1960s, *The Computer J.* 2011.

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Quantum programming primitive: Loops

[7] E. Bernstein and U. Vazirani, Quantum complexity theory, *SIAM J. on Computing* 1997

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Recursion in quantum programming

Recursive procedure in quantum programming language QPL

[8] P. Selinger, Toward a quantum programming language, *Mathematical Structures in Computer Science* 2004.

► State space of quantum system: a Hilbert space *H*

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- Quantum states: density operator an operator in \mathcal{H} : ρ is positive; $tr(\rho) = 1$.

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- Quantum states: density operator an operator in \mathcal{H} : ρ is positive; $tr(\rho) = 1$.
- Dynamics of quantum system:
 - Continuous time Schrödinger equation
 - Discrete time —
 unitary operators (closed system): UU[†] = U[†]U = I.
 super-operators (open system): Operator in the space of operators:
 completely positive; trace-preserving

Solutions to recursive equations of quantum programs Theorem:

1. The set of density operators in \mathcal{H} with the Löwner order is a CPO

[9] M. S. Ying, R. Y. Duan, Y. Feng, Z. F. Ji, Predicate transformer semantics of quantum programs, in: *Semantic Techniques in Quantum Computation*, Cambridge Univ. Press 2010.

Solutions to recursive equations of quantum programs

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 - ▶ finite-dimensional \mathcal{H} : P. Selinger (2004)

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Selinger's slogan: "Quantum data, classical control"

Control flow is classical: branching is determined by the outcomes of quantum measurements.

Example: $M = \{M_0, M_1\}$ is a quantum measurement

$$\mathbf{if} \ M[q] = 0 \to P_0$$

$$\Box \qquad 1 \to P_1$$

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"Quantum data, quantum control"

Functional quantum programming language QML, its categorical semantics

[10] T. Altenkirch and J. Grattage, A functional quantum programming language, *LICS* 2005.



How to define quantum control?

- [11] Y. Aharonov, J. Anandan, S. Popescu and L. Vaidman, Superpositions of time evolutions of a quantum system and quantum time-translation machine, *Plysical Review Letters* 1990.
- [12] A. Ambainis, E. Bach, A. Nayak, A. Vishwanath and J. Watrous, One-dimensional-quantum walks, *STOC* 2001.
- [13] D. Aharonov, A. Ambainis, J. Kempe and Vazirani, Quantum walks on graphs, *STOC* 2001.

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Introduce an external quantum coin *c*!

- state Hilbert space $\mathcal{H}_c = \text{span}\{|0\rangle, |1\rangle\}$
- ▶ U_0 and U_1 two unitary transformations on a quantum system q state Hilbert space \mathcal{H}_q .

$$\begin{aligned} \mathbf{qif} \; [c] \; |0\rangle &\rightarrow U_0[q] \\ & \qquad \Box \; |1\rangle \rightarrow U_1[q] \end{aligned}$$
 fig

Semantics of quantum case statement

▶ An unitary operator U in $\mathcal{H}_c \otimes \mathcal{H}_q$ - state Hilbert space of the composed system of coin c and principal system q:

$$U|0,\psi\rangle = |0\rangle U_0|\psi\rangle, \quad U|1,\psi\rangle = |1\rangle U_1|\psi\rangle$$

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- Quantum coin: superposition of $|0\rangle$, $|1\rangle \alpha |0\rangle + \beta |1\rangle$.
- Matrix representation:

$$U = |0\rangle\langle 0| \otimes U_0 + |1\rangle\langle 1| \otimes U_1 = \begin{pmatrix} U_0 & 0 \\ 0 & U_1 \end{pmatrix}.$$

Quantum Choice

• *W* a unitary operator in the coin's state Hilbert space \mathcal{H}_c .

[14] A. McIver and C. Morgan, Abstraction, Refinement and Proof for Probabilistic Systems, Springer 2005.

Quantum Choice

- *W* a unitary operator in the coin's state Hilbert space \mathcal{H}_c .
- ▶ Quantum choice of $U_0[q]$ and $U_1[q]$ with coin-tossing W[c]:

$$U_0[q] \oplus_{W[c]} U_1[q] \stackrel{\mathrm{def}}{=} W[c]; \ \mathbf{qif} \ [c] \ |0\rangle \to U_0[q]$$

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Compare with probabilistic choice!

[14] A. McIver and C. Morgan, Abstraction, Refinement and Proof for Probabilistic Systems, Springer 2005.



A more general quantum case statement

$$\begin{array}{c} \mathbf{qif} \ [c] \ |0\rangle \rightarrow P_0 \\ & \Box \ |1\rangle \rightarrow P_1 \end{array}$$
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 $ightharpoonup P_0$, P_1 include quantum measurements.

[15] Chapter 6 of M. S. Ying, Foundations of Quantum Programming, Elsevier - Morgan Kaufmann 2016.

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- How to define the semantics?

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One-dimensional quantum walk

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 - $\mathcal{H}_d = \text{span}\{|L\rangle, |R\rangle\}$, L, R indicate the direction Left and Right.
 - $\mathcal{H}_p = \text{span}\{|n\rangle : n \in \mathbb{Z}\}$, n indicates the position marked by integer n.

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► Hadamard walk —— repeated applications of operator *W*.

One-dimensional quantum walk — a different view

▶ Define the left and right translation operators T_L and T_R in the position space \mathcal{H}_p :

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One-dimensional quantum walk — a different view

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$$T = \mathbf{qif} [d] |L\rangle \rightarrow T_L[p]$$

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$$\mathbf{fiq}$$

► The single-step walk operator *U* is a quantum choice:

$$T_L[p] \oplus_{H[d]} T_R[p]$$

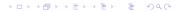
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- ► The walk —— a recursive program *X* declared by the recursive equation:

$$X \Leftarrow T_L[p] \oplus_{H[d]} (T_R[p]; X)$$



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 - if d is in state |R⟩ then it moves one position right, also followed by
 a procedure behaving as the recursive walk itself.
- ► The walk a program *X* (or *Y*) declared by the equation:

$$X \Leftarrow (T_L[p]; X) \oplus_{H[d]} (T_R[p]; X)$$

A More Interesting Recursive Quantum Walk

Let $n \ge 2$. A variant of unidirectional recursive quantum walk:

$$X \leftarrow ((T_L[p];X) \oplus_{H[d]} (T_R[p];X)); (T_L[p] \oplus_{H[d]} T_R[p])^n$$

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How to solve these quantum recursive equations?

Syntactic Approximation

► A recursive program *X* declared by equation

$$X \Leftarrow F(X)$$

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Syntactic approximations:

$$\begin{cases} X^{(0)} = \textbf{Abort,} \\ X^{(n+1)} = F[X^{(n)}/X] \text{ for } n \ge 0. \end{cases}$$

Program $X^{(n)}$ is the nth syntactic approximation of X.

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Program $X^{(n)}$ is the *n*th syntactic approximation of X.

• Semantics [X] of X is the limit

$$\llbracket X \rrbracket = \lim_{n \to \infty} \llbracket X^{(n)} \rrbracket$$

```
\begin{split} X^{(0)} &= \textbf{abort}, \\ X^{(1)} &= T_L[p] \oplus_{H[d]} (T_R[p]; \textbf{abort}), \\ X^{(2)} &= T_L[p] \oplus_{H[d]} (T_R[p]; T_L[p] \oplus_{H[d_1]} (T_R[p]; \textbf{abort})), \\ X^{(3)} &= T_L[p] \oplus_{H[d]} (T_R[p]; T_L[p] \oplus_{H[d_1]} (T_R[p]; T_L[p] \oplus_{H[d_2]} (T_R[p]; \textbf{abort}))), \end{split}
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Observations

- Continuously introduce new coin to avoid variable conflict.
- ▶ Variables d, d₁, d₂, ... denote identical particles.

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- ► The number of the coin particles that are needed in running the recursive walk is unknown beforehand.

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Observations

- Continuously introduce new coin to avoid variable conflict.
- ▶ Variables d, d₁, d₂, ... denote identical particles.
- The number of the coin particles that are needed in running the recursive walk is unknown beforehand.
- ► We need to deal with *quantum systems* where the number of particles of the same type may vary.

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- Let \mathcal{H} be the state Hilbert space of one particle.
- ► For each permutation π of 1, ..., n, define the permutation operator P_{π} in $\mathcal{H}^{\otimes n}$:

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$$P_{\pi}|\psi_1\otimes...\otimes\psi_n\rangle=|\psi_{\pi(1)}\otimes...\otimes\psi_{\pi(n)}\rangle$$

• Define the symmetrisation and antisymmetrisation operators in $\mathcal{H}^{\otimes n}$:

$$S_{+} = \frac{1}{n!} \sum_{\pi} P_{\pi}, \quad S_{-} = \frac{1}{n!} \sum_{\pi} (-1)^{\pi} P_{\pi}$$

- v = + for bosons, v = for fermions.
 - ► Symmetrisation or antisymmetrisation:

$$|\psi_1,...,\psi_n\rangle_v=S_v|\psi_1\otimes...\otimes\psi_n\rangle.$$

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▶ State space of *n* bosons and that of fermions:

$$\mathcal{H}_v^{\otimes n} = S_v \mathcal{H}^{\otimes n} = \operatorname{span}\{|\psi_1,...,\psi_n\rangle_v : |\psi_1\rangle,...,|\psi_n\rangle \text{ are in } \mathcal{H}\}$$

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► Introduce the vacuum state |**0**⟩:

$$\mathcal{H}_v^{\otimes 0} = \mathcal{H}^{\otimes 0} = \operatorname{span}\{|\mathbf{0}\rangle\}.$$

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Introduce the vacuum state |0>:

$$\mathcal{H}_v^{\otimes 0} = \mathcal{H}^{\otimes 0} = \text{span}\{|\mathbf{0}\rangle\}.$$

► The space of the states of variable particle number is the Fock space:

$$\mathcal{F}_v(\mathcal{H}) = \sum_{n=0}^{\infty} \mathcal{H}_v^{\otimes n}$$

ightharpoonup (discrete-time) evolution of one particle —— unitary operator U.

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▶ Extend to the Fock spaces $\mathcal{F}_v(\mathcal{H})$ and $\mathcal{F}(\mathcal{H})$:

$$\mathbf{U}\left(\sum_{n=0}^{\infty}|\Psi(n)\rangle\right)=\sum_{n=0}^{\infty}\mathbf{U}|\Psi(n)\rangle$$

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► Annihilation operator $a(\psi)$ — the Hermitian conjugate of $a^*(\psi)$:

$$a(\psi)|\mathbf{0}\rangle = 0,$$

$$a(\psi)|\psi_1,...,\psi_n\rangle_v = \frac{1}{\sqrt{n}} \sum_{i=1}^n (v)^{i-1} \langle \psi|\psi_i\rangle |\psi_1,...,\psi_{i-1},\psi_{i+1},...,\psi_n\rangle_v$$

Decrease the number of particles by one unit, while preserving the symmetry of the state.

Outline

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- 2. Recursive Programming
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Example - Unidirectional Recursive Hadamard Walk

Semantics of the recursive Hadamard walk:

$$\llbracket X \rrbracket = \left[\sum_{i=0}^{\infty} \left(\bigotimes_{j=0}^{i-1} |R\rangle_{d_j} \langle R| \otimes |L\rangle_{d_i} \langle L| \right) \otimes T_L T_R^i \right] (\mathbf{H} \otimes I)$$

An operator in

$$\mathcal{F}_v(\mathcal{H}_d) \otimes \mathcal{H}_p \to \mathcal{F}(\mathcal{H}_d) \otimes \mathcal{H}_p.$$

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► The sign v is + or -, depending on using bosons or fermions to implement the direction coins d, d₁, d₂,

Principal System Semantics

• Each state $|\Psi\rangle$ in Fock space $\mathcal{F}_v(\mathcal{H}_d)$ induces mapping:

$$[\![X, \Psi]\!]_p$$
: pure states \to partial density operators in \mathcal{H}_p
 $[\![X, \Psi]\!]_p(|\psi\rangle) = tr_{\mathcal{F}(\mathcal{H}_d)}(|\Phi\rangle\langle\Phi|)$

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▶ $[\![X, \Psi]\!]_p$ is called the principal system semantics of X with coin initialisation $|\Psi\rangle$.

Example - Bidirectional Recursive Quantum Walk

$$\begin{cases} X \Leftarrow T_L[p] \oplus_{H[d]} (T_R[p]; Y), \\ Y \Leftarrow (T_L[p]; X) \oplus_{H[d]} T_R[p] \end{cases}$$

► Coherent state of bosons in the symmetric Fock space $\mathcal{F}_+(\mathcal{H})$ over \mathcal{H} :

$$|\psi\rangle_{\mathrm{coh}} = \exp\left(-\frac{1}{2}\langle\psi|\psi\rangle\right) \sum_{n=0}^{\infty} \frac{[a^*(\psi)]^n}{n!} |\mathbf{0}\rangle$$

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➤ The walk starts from position 0 and the coins are initialised in the coherent states of bosons corresponding to |L⟩:

$$\begin{split} \llbracket X, L_{\mathrm{coh}} \rrbracket_p(|0\rangle) &= \frac{1}{\sqrt{e}} \left(\sum_{k=0}^{\infty} \frac{1}{2^{2k+1}} |-1\rangle \langle -1| + \sum_{k=0}^{\infty} \frac{1}{2^{2k+2}} |2\rangle \langle 2| \right) \\ &= \frac{1}{\sqrt{e}} \left(\frac{2}{3} |-1\rangle \langle -1| + \frac{1}{3} |2\rangle \langle 2| \right). \end{split}$$

[16] Chapter 7 of M. S. Ying, Foundations of Quantum Programming,

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Quantum programming theory

- ► Imperative quantum programs
 - [1] M. S. Ying, Floyd-Hoare logic for quantum programs, *TOPLAS* 2011.
 - [2] M. S. Ying, S. G. Ying and X. D. Wu, Invariants of quantum programs: characterisations and generation, *POPL* 2017.

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- ► Functional quantum programs
 - [1] M. Pagani, P. Selinger and B. Valiron, Applying quantitative semantics to higher-order quantum computing, *POPL* 2014.
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THANK YOU!